

Towards a combinatorial characterization of equistable graphs

(Partial results on a conjecture of Orlin)

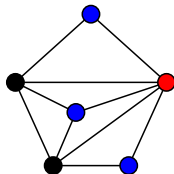
Martin Milanič

UP FAMNIT and UP PINT, University of Primorska

Raziskovalni matematični seminar, FAMNIT, 28. marec 2011

Graphs and stable sets

- $G = (V, E)$ - a finite simple undirected graph
- **stable (independent) set**: a subset $S \subseteq V$ of pairwise non-adjacent vertices
- a stable set is **maximal** if it is not contained in any larger stable set



Definition

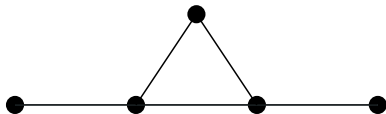
A graph $G = (V, E)$ is **equistable** if there exists a function $w : V \rightarrow \mathbb{N}$ and a positive integer t such that

$\forall S \subseteq V$:

S is a maximal stable set in $G \iff w(S) = \sum_{v \in S} w(v) = t.$

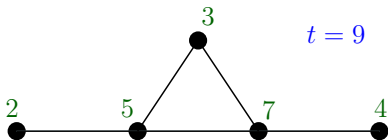
Equistable graphs: example

The following graph is equistable:



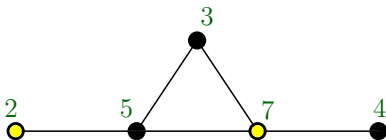
Equistable graphs: example

The following graph is equistable:



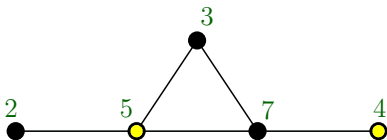
Equistable graphs: example

The following graph is equistable:



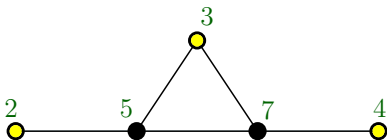
Equistable graphs: example

The following graph is equistable:



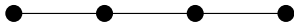
Equistable graphs: example

The following graph is equistable:



Equistable graphs: example

The following graph is not equistable:



Equistable graphs: example

The following graph is not equistable:



If

$$w_1 + w_3 = t$$

$$w_2 + w_4 = t$$

$$w_1 + w_4 = t$$

then

$$w_2 + w_3 = t.$$

Equistable graphs: example

The following graph is not equistable:



If

$$w_1 + w_3 = t$$

$$w_2 + w_4 = t$$

$$w_1 + w_4 = t$$

then

$$w_2 + w_3 = t.$$

Equistable graphs: motivation

- **threshold graphs** (Chvátal-Hammer 1977):
 $\exists w, t$ s.t. $S \subseteq V$ **stable** $\Leftrightarrow w(S) \leq t$
- **equistable graphs** (Payan 1980):
 $\exists w, t$ s.t. $S \subseteq V$ **maximal stable** $\Leftrightarrow w(S) = t$

Equistable graphs generalize:

- threshold graphs (Payan, 1980);
- co-graphs (graphs without an induced 3-edge path) (Mahadev-Peled-Sun, 1994).

Equistable graphs: motivation

- **threshold graphs** (Chvátal-Hammer 1977):
 $\exists w, t$ s.t. $S \subseteq V$ **stable** $\Leftrightarrow w(S) \leq t$
- **equistable graphs** (Payan 1980):
 $\exists w, t$ s.t. $S \subseteq V$ **maximal stable** $\Leftrightarrow w(S) = t$

Equistable graphs generalize:

- threshold graphs (Payan, 1980);
- co-graphs (graphs without an induced 3-edge path) (Mahadev-Peled-Sun, 1994).

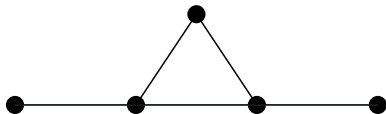
Definition

A graph $G = (V, E)$ is a **general partition graph (gpg)** if there exists a finite set U and an assignment of nonempty subsets $U_x \subseteq U$ to vertices of V such that

- $xy \in E$ if and only if $U_x \cap U_y \neq \emptyset$, and
- for every maximal stable set S in G , the set $\{U_x : x \in S\}$ forms a partition of U .

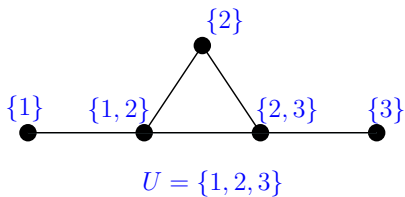
General partition graphs: example

The following graph is a gpg:



General partition graphs: example

The following graph is a gpg:



General partition graphs: a characterization

Theorem (McAvaney-Robertson-DeTemple, 1993)

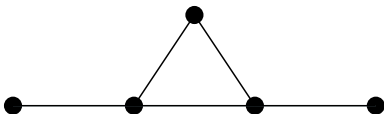
For every graph G , the following are equivalent:

- *G is a gpg,*
- *every edge of G is contained in a strong clique.*

strong clique = a clique meeting all maximal stable sets

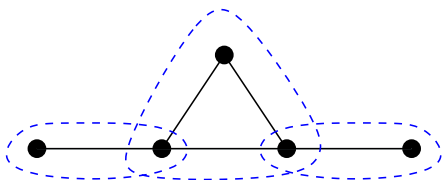
General partition graphs: example

The following graph is a gpg (every edge is contained in a strong clique):



General partition graphs: example

The following graph is a gpg (every edge is contained in a strong clique):



General partition graphs: example

The following graph is not a gpg (there exists an edge not contained in any strong clique):



General partition graphs: example

The following graph is not a gpg (there exists an edge not contained in any strong clique):



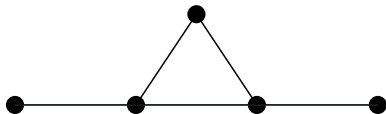
Definition

A graph $G = (V, E)$ is a **triangle graph** if it satisfies the following triangle condition:

- for every maximal stable set S in G and every edge $uv \in E(G - S)$, u and v have a common neighbor in S .

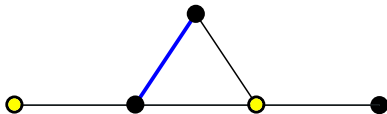
Triangle graphs: example

The following graph is a triangle graph:



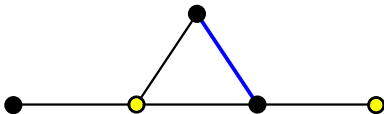
Triangle graphs: example

The following graph is a triangle graph:



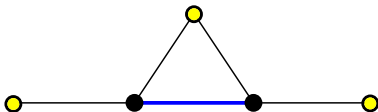
Triangle graphs: example

The following graph is a triangle graph:



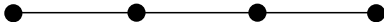
Triangle graphs: example

The following graph is a triangle graph:



Triangle graphs: example

The following graph is not a triangle graph:



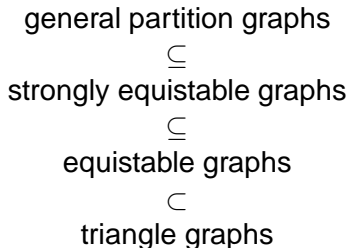
Triangle graphs: example

The following graph is not a triangle graph:



Inclusion relations among these classes

The following inclusion relations hold:



A condition equivalent to the triangle condition

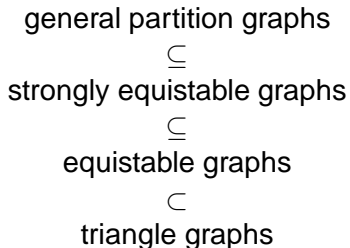
In the equistable graphs literature, the triangle condition was replaced with the following equivalent condition:

“absence of a **bad** P_4 ”:

For each induced P_4 on the vertices a, b, c, d , each maximal stable set containing the end-vertices a and d has a common neighbor of the middle vertices b and c .

Inclusion relations among these classes

The following inclusion relations hold:



General partition graphs are equistable

Theorem (Jim Orlin, 2009)

Every gpg is equistable.

Proof idea:

Let $(U_x : x \in V)$ with $U = \{u_1, \dots, u_k\}$ be a set system realizing G .

Let $n = |V(G)|$ and assign to each $x \in V$ weight

$$w(x) = \sum \{n^j : u_j \in U_x\},$$

also let

$$t = \sum_{j=1}^k n^j.$$

General partition graphs are equistable

Theorem (Jim Orlin, 2009)

Every gpg is equistable.

Proof idea:

Let $(U_x : x \in V)$ with $U = \{u_1, \dots, u_k\}$ be a set system realizing G .

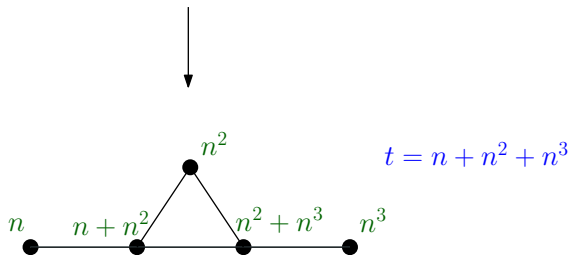
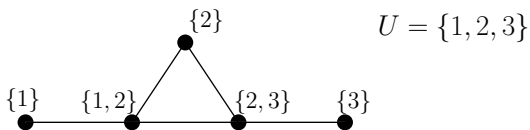
Let $n = |V(G)|$ and assign to each $x \in V$ weight

$$w(x) = \sum \{n^j : u_j \in U_x\},$$

also let

$$t = \sum_{j=1}^k n^j.$$

Example



General partition graphs are strongly equitable

Theorem (Jim Orlin, 2009)

Every gpg is equitable.

Theorem (McAvaney-Robertson-DeTemple, 1993)

G is a gpg if and only if every edge of G is contained in a strong clique.

Theorem (Mahadev-Peled-Sun, 1994)

Every equitable with a strong clique is strongly equitable.

Corollary

Every gpg is strongly equitable.

General partition graphs are strongly equitable

Theorem (Jim Orlin, 2009)

Every gpg is equitable.

Theorem (McAvaney-Robertson-DeTemple, 1993)

G is a gpg if and only if every edge of G is contained in a strong clique.

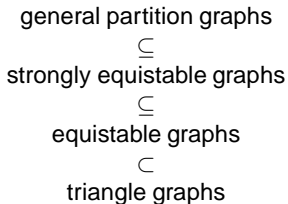
Theorem (Mahadev-Peled-Sun, 1994)

Every equitable with a strong clique is strongly equitable.

Corollary

Every gpg is strongly equitable.

Inclusion relations among these classes



Conjecture (Mahadev-Peled-Sun, 1994)

Every equistable graph is strongly equistable.

Conjecture (Jim Orlin, 2009)

Every equistable graph is a gpg.

If true, Orlin's conjecture would provide a combinatorial characterization of equistable graphs.

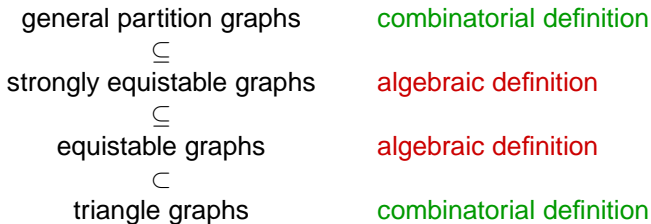
Orlin's conjecture holds within the following graph classes:

- chordal graphs,
- graphs obtained from triangle-free graphs by gluing chordal graphs along edges,
- outerplanar graphs,
- series-parallel graphs $\equiv K_4$ -minor-free graphs.

Orlin's conjecture holds within the following graph classes:

- chordal graphs,
- graphs obtained from triangle-free graphs by gluing chordal graphs along edges,
- outerplanar graphs,
- series-parallel graphs $\equiv K_4$ -minor-free graphs.

Inclusion relations among these classes

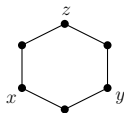
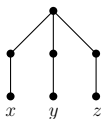
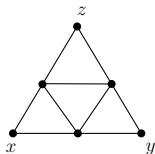


Our goal:

Identify further combinatorially defined graph classes \mathcal{C} such that within \mathcal{C} , some of the above inclusions become equalities.

First example: AT-free graphs

Asteroidal triple (AT): a triple of vertices x, y, z such that for every pair of them there is a path connecting the two vertices that avoids the neighborhood of the remaining vertex.



Theorem (Kloks-Lee-Liu-Müller, 2003)

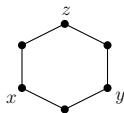
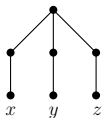
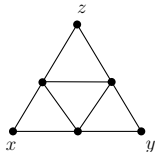
Every AT-free triangle graph is a gpg.

Corollary

Orlin's conjecture holds within the class of AT-free graphs.

First example: AT-free graphs

Asteroidal triple (AT): a triple of vertices x, y, z such that for every pair of them there is a path connecting the two vertices that avoids the neighborhood of the remaining vertex.



Theorem (Kloks-Lee-Liu-Müller, 2003)

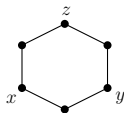
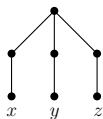
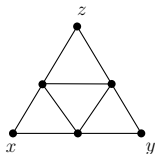
Every AT-free triangle graph is a gpg.

Corollary

Orlin's conjecture holds within the class of AT-free graphs.

First example: AT-free graphs

Asteroidal triple (AT): a triple of vertices x, y, z such that for every pair of them there is a path connecting the two vertices that avoids the neighborhood of the remaining vertex.



Theorem (Kloks-Lee-Liu-Müller, 2003)

Every AT-free triangle graph is a gpg.

Corollary

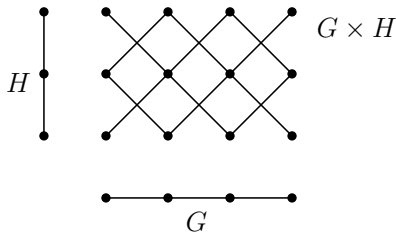
Orlin's conjecture holds within the class of AT-free graphs.

Tensor products

Definition

The **tensor** (or **direct**) **product** of graphs G and H is the graph $G \times H$ such that:

- $V(G \times H) = V(G) \times V(H)$,
- $(u, x)(v, y) \in E(G \times H)$ if and only if $uv \in E(G) \wedge xy \in E(H)$.



Theorem

Let $G = G_1 \times G_2$, where G_1, G_2 are connected, with more than just one vertex. The following are equivalent:

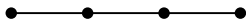
- (i) G is a gpg.*
- (ii) G is strongly equistable.*
- (iii) G is equistable.*
- (iv) $\exists m \geq 2$ such that G_1 and G_2 are complete m -partite graphs.*

Proof sketch of (iii) \Rightarrow (iv)

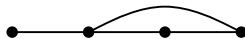
$G = G_1 \times G_2$ equistable.

Claim

No induced subgraph of G_1 is isomorphic to P_4 or a paw.



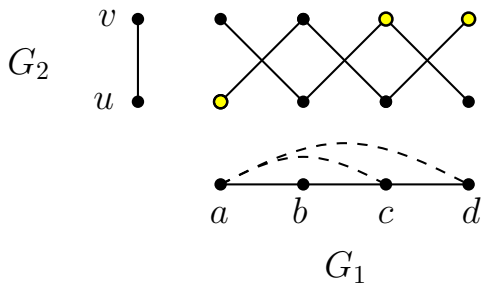
P_4



paw

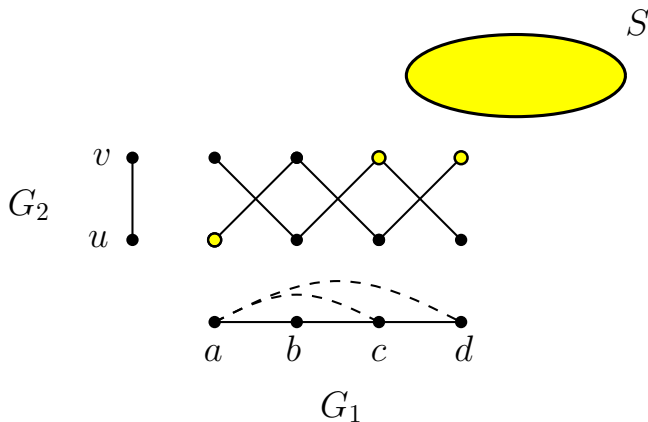
Proof sketch of (iii) \Rightarrow (iv)

$G = G_1 \times G_2$ equistable.



Proof sketch of (iii) \Rightarrow (iv)

S : a maximal stable set of G containing (a, u) , (c, v) , (d, v)



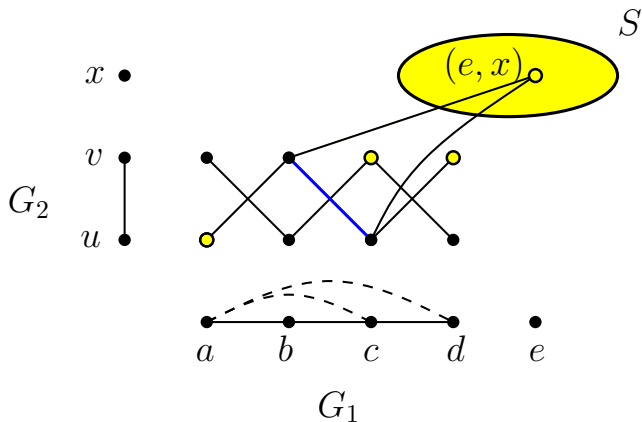
Proof sketch of (iii) \Rightarrow (iv)

Every equistable graph satisfies the triangle condition:

- for every maximal stable set S in G and every edge $uv \in E(G - S)$, u and v have a common neighbor in S .

Proof sketch of (iii) \Rightarrow (iv)

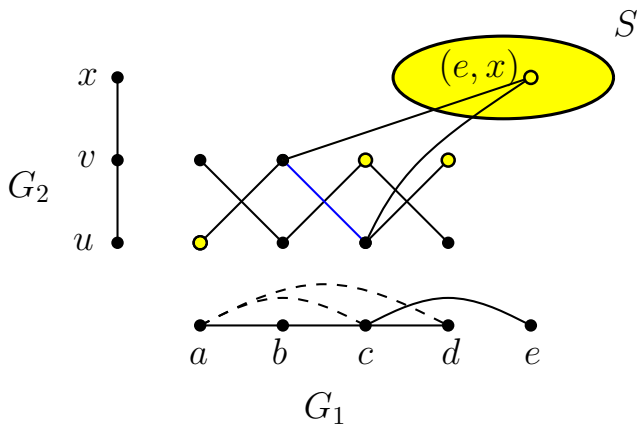
$(e, x) \in S$: a common neighbor of (b, v) and (c, u)



Proof sketch of (iii) \Rightarrow (iv)

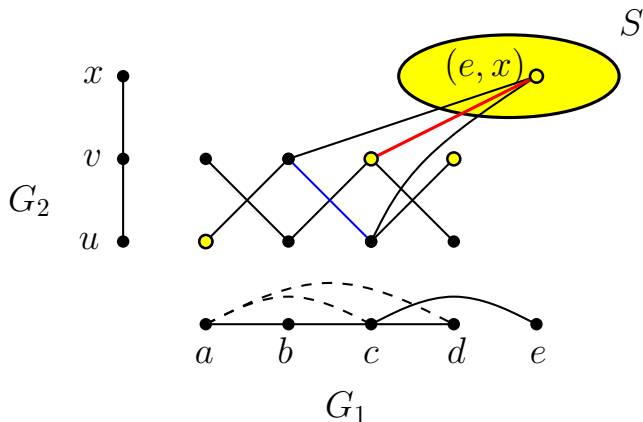
$$(e, x)(c, u) \in E(G_1 \times G_2) \Rightarrow ce \in E(G_1)$$

$$(e, x)(b, v) \in E(G_1 \times G_2) \Rightarrow vx \in E(G_2)$$



Proof sketch of (iii) \Rightarrow (iv)

$$ce \in E(G_1) \wedge vx \in E(G_2) \Rightarrow (c, v)(e, x) \in E(G_1 \times G_2)$$



Contradiction with the fact that S is a stable set.

Proof sketch of (iii) \Rightarrow (iv)

$G = G_1 \times G_2$ equistable.

Claim

G_1 is (P_4, paw) -free.

G_2 is (P_4, paw) -free.

Lemma

Every connected (P_4, paw) -free graph is complete multipartite.

So both G_1 and G_2 are complete multipartite.

Proof sketch of (iii) \Rightarrow (iv)

$G = G_1 \times G_2$ equistable.

Claim

G_1 is (P_4, paw) -free.

G_2 is (P_4, paw) -free.

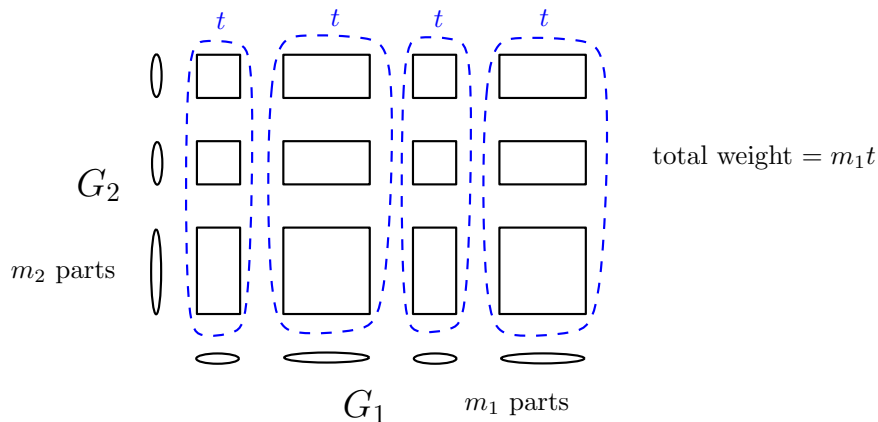
Lemma

Every connected (P_4, paw) -free graph is complete multipartite.

So both G_1 and G_2 are complete multipartite.

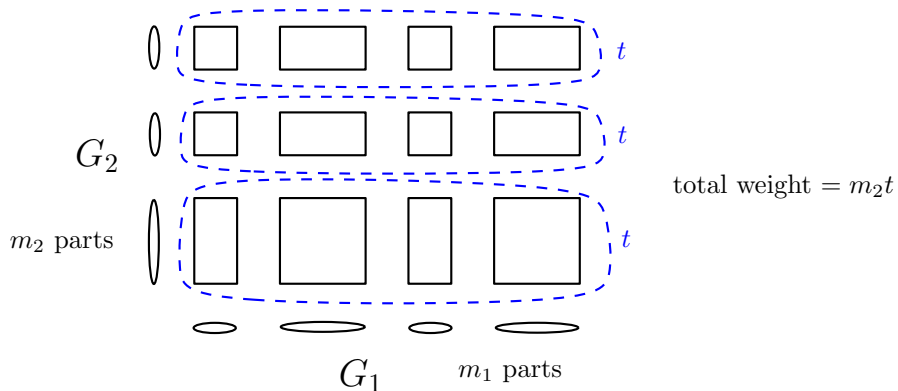
Proof sketch of (iii) \Rightarrow (iv)

G_1 and G_2 have the same number of parts:



Proof sketch of (iii) \Rightarrow (iv)

G_1 and G_2 have the same number of parts:

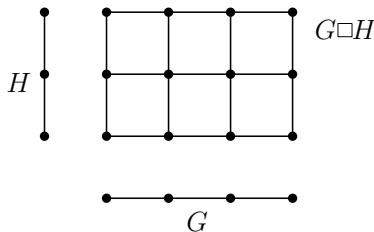


Cartesian products

Definition

The **Cartesian product** of graphs G and H is the graph $G \square H$ such that:

- $V(G \square H) = V(G) \times V(H)$,
- $(u, x)(v, y) \in E(G \square H)$ if and only if $(u = v \wedge xy \in E(H)) \vee (x = y \wedge uv \in E(G))$.



Theorem

Let $G = G_1 \square G_2$, where G_1, G_2 are connected, with more than just one vertex. The following are equivalent:

- (i) G is a *gpg*.
- (ii) G is strongly equistable.
- (iii) G is equistable.
- (iv) G is a triangle graph.
- (v) $\exists m \geq 2$ such that $G_1 \cong G_2 \cong K_m$.

Proof sketch of $(iv) \Rightarrow (v)$

For a graph K , we denote:

- by $\delta(K)$ the minimum vertex degree,
- by $\Delta(K)$ the maximum vertex degree,
- for $uv \in E(K)$, let $\lambda(uv) = |N(u) \cap N(v)|$,
- $\lambda(K) = \min\{\lambda(uv) : uv \in E(K)\}$.

Observation

$$\lambda(K) \leq \delta(K) - 1.$$

Proof sketch of $(iv) \Rightarrow (v)$

$G = G_1 \square G_2$ triangle graph

Claim

$$\lambda(G_2) \geq \Delta(G_1) - 1.$$

Suppose for a contradiction that $\lambda(G_2) \leq \Delta(G_1) - 2$.

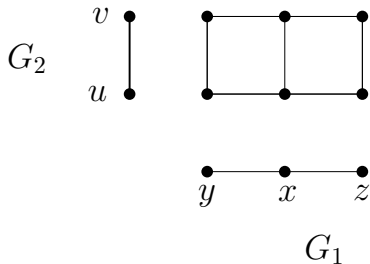
Proof sketch of $(iv) \Rightarrow (v)$

$$\lambda(G_2) \leq \Delta(G_1) - 2.$$

x : vertex of maximum degree in G_1 ($d(x) \geq 2$)

y, z : two neighbors of x ,

$uv \in E(G_2)$: $\lambda(uv) = \lambda(G_2)$.



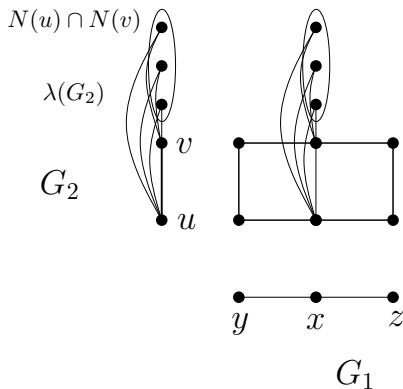
Proof sketch of $(iv) \Rightarrow (v)$

$$\lambda(G_2) \leq \Delta(G_1) - 2.$$

x : vertex of maximum degree in G_1 ($d(x) \geq 2$)

y, z : two neighbors of x ,

$uv \in E(G_2)$: $\lambda(uv) = \lambda(G_2)$.



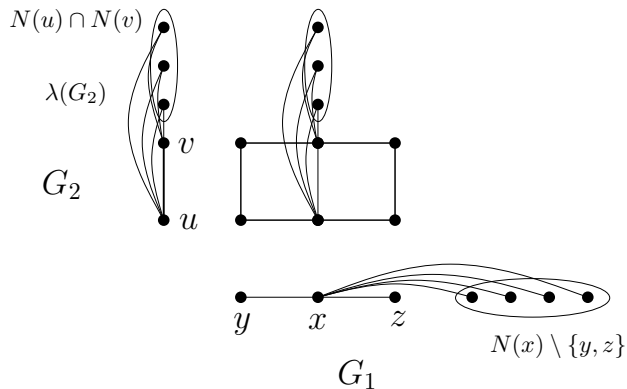
Proof sketch of $(iv) \Rightarrow (v)$

$$\lambda(G_2) \leq \Delta(G_1) - 2.$$

x : vertex of maximum degree in G_1 ($d(x) \geq 2$)

y, z : two neighbors of x ,

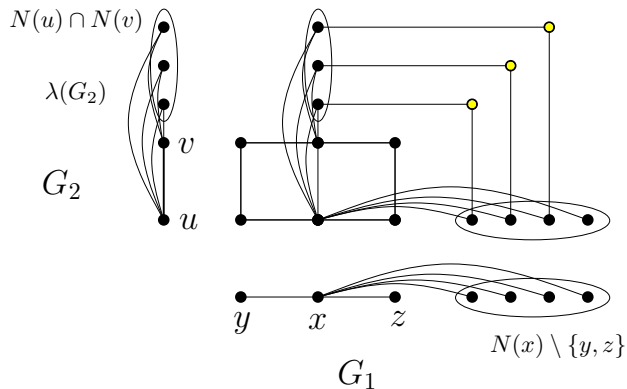
$uv \in E(G_2)$: $\lambda(uv) = \lambda(G_2)$.



Proof sketch of $(iv) \Rightarrow (v)$

$$\lambda(G_2) \leq \Delta(G_1) - 2.$$

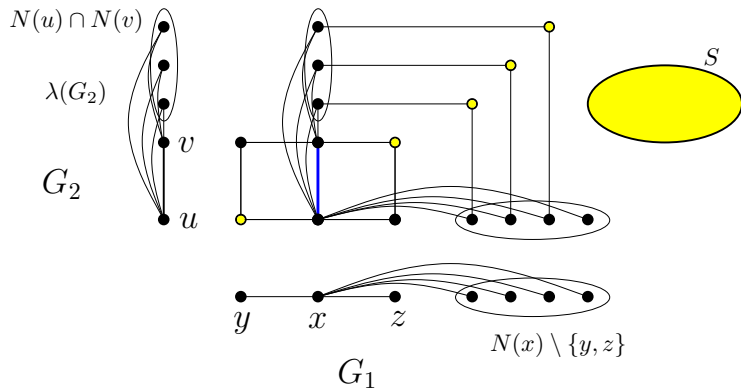
$$|N(x) \setminus \{y, z\}| \geq \lambda(G_2).$$



Proof sketch of $(iv) \Rightarrow (v)$

$$\lambda(G_2) \leq \Delta(G_1) - 2.$$

S : a maximal stable set containing all yellow vertices.



Contradiction.

Proof sketch of $(iv) \Rightarrow (v)$

Claim

$$\lambda(G_2) \geq \Delta(G_1) - 1.$$

$$\lambda(G_1) \geq \Delta(G_2) - 1.$$

Consequently

$$\lambda(G_1) \geq \Delta(G_2) - 1 \geq \delta(G_2) - 1 \geq \lambda(G_2) \geq \Delta(G_1) - 1 \geq \delta(G_1) - 1 \geq \lambda(G_1).$$

Equalities hold, therefore:

- $\delta(G_1) = \Delta(G_1) = \delta(G_2) = \Delta(G_2)$.
- G_1 and G_2 are both regular of the same degree $m - 1$.
- $\lambda(G_i) = m - 2$ implies that $G_1 \cong G_2 \cong K_m$.

Proof sketch of $(iv) \Rightarrow (v)$

Claim

$$\lambda(G_2) \geq \Delta(G_1) - 1.$$

$$\lambda(G_1) \geq \Delta(G_2) - 1.$$

Consequently

$$\lambda(G_1) \geq \Delta(G_2) - 1 \geq \delta(G_2) - 1 \geq \lambda(G_2) \geq \Delta(G_1) - 1 \geq \delta(G_1) - 1 \geq \lambda(G_1).$$

Equalities hold, therefore:

- $\delta(G_1) = \Delta(G_1) = \delta(G_2) = \Delta(G_2)$.
- G_1 and G_2 are both regular of the same degree $m - 1$.
- $\lambda(G_i) = m - 2$ implies that $G_1 \cong G_2 \cong K_m$.

Proof sketch of $(iv) \Rightarrow (v)$

Claim

$$\lambda(G_2) \geq \Delta(G_1) - 1.$$

$$\lambda(G_1) \geq \Delta(G_2) - 1.$$

Consequently

$$\lambda(G_1) \geq \Delta(G_2) - 1 \geq \delta(G_2) - 1 \geq \lambda(G_2) \geq \Delta(G_1) - 1 \geq \delta(G_1) - 1 \geq \lambda(G_1).$$

Equalities hold, therefore:

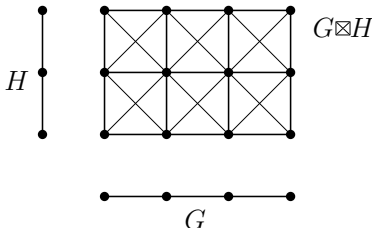
- $\delta(G_1) = \Delta(G_1) = \delta(G_2) = \Delta(G_2)$.
- G_1 and G_2 are both regular of the same degree $m - 1$.
- $\lambda(G_i) = m - 2$ implies that $G_1 \cong G_2 \cong K_m$.

Strong products

Definition

The **strong product** of graphs G and H is the graph $G \boxtimes H$ such that:

- $V(G \boxtimes H) = V(G) \times V(H)$,
- $E(G \boxtimes H) = E(G \times H) \cup E(G \square H)$.



Theorem (McAvaney-Robertson-DeTemple, 1993)

If $G_1 \boxtimes G_2$ is a gpg, then G_1 and G_2 are gpgs.

Proposition

The general partition (equistable, strongly equistable, triangle) graphs are not closed under the strong product.

Counterexample: $(K_3 \square K_3) \boxtimes (K_3 \square K_3)$.

Theorem (McAvaney-Robertson-DeTemple, 1993)

If $G_1 \boxtimes G_2$ is a gpg, then G_1 and G_2 are gpgs.

Proposition

The general partition (equistable, strongly equistable, triangle) graphs are not closed under the strong product.

Counterexample: $(K_3 \square K_3) \boxtimes (K_3 \square K_3)$.

Equistable strong products

Theorem

If $G_1 \boxtimes G_2$ is a triangle graph, then G_1 and G_2 are triangle graphs.

Theorem

If every edge of G_1 and G_2 is contained in a simplicial clique, then $G_1 \boxtimes G_2$ is a gpg.

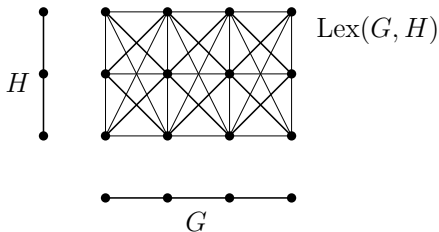
simplicial clique: a clique of the form $K = N(v) \cup \{v\}$.

Lexicographic products

Definition

The **lexicographic product** of graphs G and H is the graph $\text{Lex}(G, H)$ such that:

- $V(\text{Lex}(G, H)) = V(G) \times V(H)$,
- $(u, x)(v, y) \in E(\text{Lex}(G, H))$ if and only if $(uv \in E(G)) \vee (u = v \wedge xy \in E(H))$.



Theorem (McAvaney-Robertson-DeTemple, 1993)

$\text{Lex}(G_1, G_2)$ is a *gpg* if and only if G_1 and G_2 are *gpgs*.

Theorem

$\text{Lex}(G_1, G_2)$ is a *triangle graph* if and only if G_1 and G_2 are *triangle graphs*.

Theorem

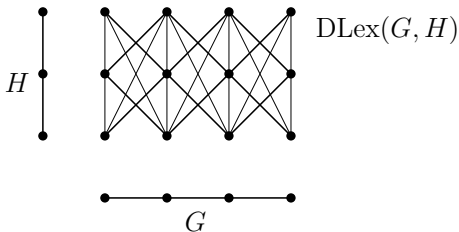
- (i) If $\text{Lex}(G_1, G_2)$ is *equistable*, then G_1 and G_2 are *equistable*.
- (ii) If G_1 and G_2 are *equistable* and G_2 contains an *isolated vertex*, then $\text{Lex}(G_1, G_2)$ is *equistable*.

Deleted lexicographic products

Definition

The **deleted lexicographic product** of graphs G and H is the graph $\text{DLex}(G, H)$ such that:

- $V(\text{DLex}(G, H)) = V(G) \times V(H)$,
- $(u, x)(v, y) \in E(\text{DLex}(G, H))$ if and only if $(uv \in E(G) \wedge x \neq y) \vee (u = v \wedge xy \in E(H))$.



Theorem

Let G be connected, triangle-free, with at least two vertices, H with at least one edge. The following are equivalent:

- (i) $\text{DLex}(G, H)$ is a gpg.
- (ii) $\text{DLex}(G, H)$ is strongly equistable.
- (iii) $\text{DLex}(G, H)$ is equistable.
- (iv) $\text{DLex}(G, H)$ is a triangle graph.
- (v) Either $G = K_2$ and the complement of H is of maximum degree at most 1, or $G \neq K_2$ is complete bipartite and $H = K_{t \times 2}$ for some $t \geq 2$.

Open problems

Determine the complexity of recognizing equistable graphs.

Conjecture (Mahadev-Peled-Sun, 1994)

Every equistable graph is strongly equistable.

Conjecture (Jim Orlin, 2009)

Every equistable graph is a gpg.

Conjecture (ŠM-MM, 2011)

Every equistable graph contains a strong clique.

Thank you