

# Equistarable Bipartite Graphs

Nina Chiarelli

Joint work with Endre Boros and Martin Milanič

FAMNIT, Koper, November 2014

# Outline

- 1 Equistable graphs
- 2 Equistarable graphs
- 3 Special cases
  - Bipartite graphs
  - Forests

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# Equistable graphs

Payan, 1980

A **stable set** (or independent set) in a graph  $G = (V, E)$  is a subset  $S \subseteq V$  such that **no two vertices in  $S$  are adjacent**.

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## Definition

A graph  $G$  is said to be **equistable** if there exists a mapping  $\varphi : V \rightarrow [0, 1]$  such that for all  $S \subseteq V$ ,

$$S \text{ is a maximal stable set} \iff \varphi(S) := \sum_{v \in S} \varphi(v) = 1.$$

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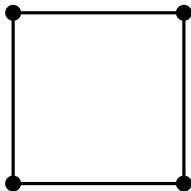
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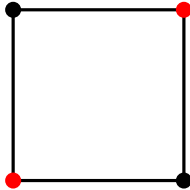
Such a  $\varphi$  is called an **equistable weight function** of  $G$ .

## Example of an equistable graph

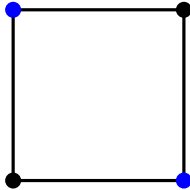




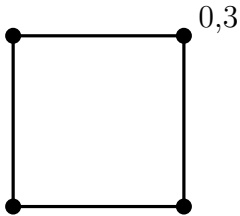
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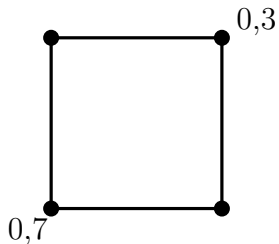
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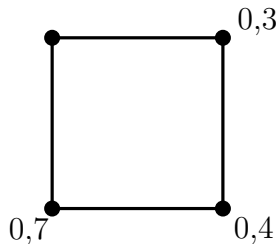
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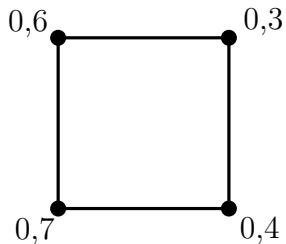
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## Example of an equistable graph



# Equistable graphs

in connection to some other graph classes

Equistable

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in connection to some other graph classes

Mahadev et al.

Strongly equistable

Equistable



# Equistable graphs

in connection to some other graph classes

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## Definition

Given a graph  $G$ , let  $\mathcal{S}(G)$  be the set of all maximal stable sets of  $G$ , and  $\mathcal{T}(G)$  the set of all other nonempty subsets of  $V(G)$ .

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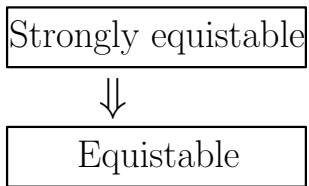
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General partition

Strongly equistable



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DeTemple et al.

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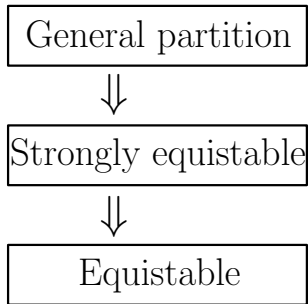
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Theorem (McAvaney et al.,  
1993)

A graph  $G$  is a **general partition graph** if and only if every edge of  $G$  is contained in a strong clique.

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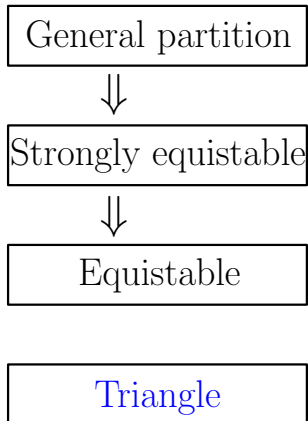
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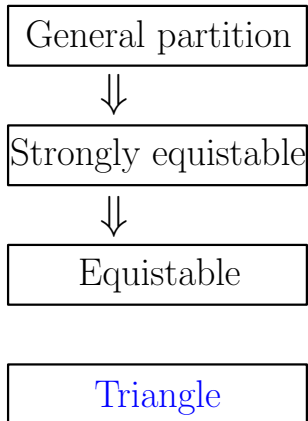
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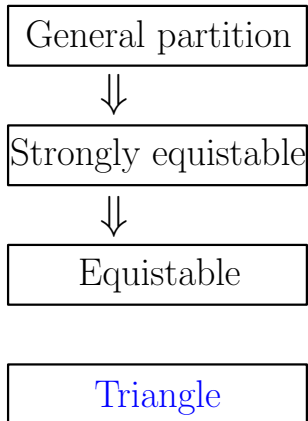
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## Triangle condition

For every maximal stable set  $S$  in  $G = (V, E)$  and every edge  $uv$  in  $G - S$  there is a vertex  $s \in S$  such that  $\{u, v, s\}$  induces a triangle in  $G$ .

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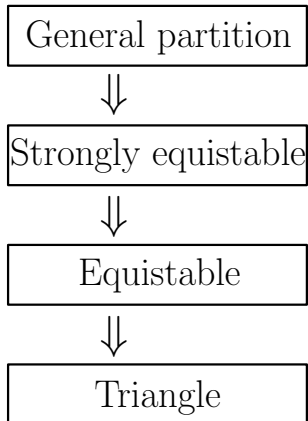
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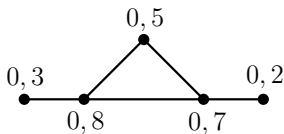
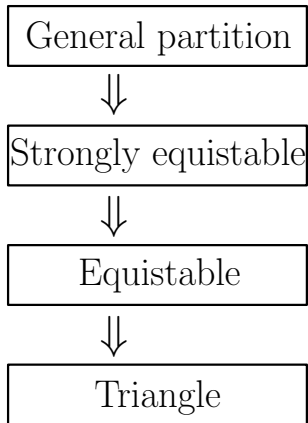
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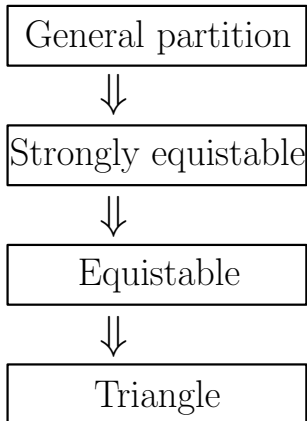
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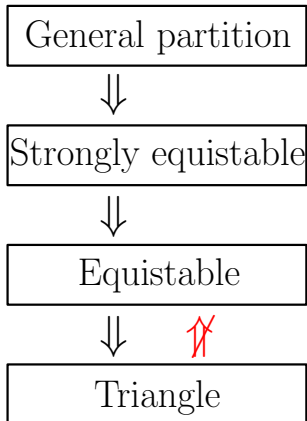
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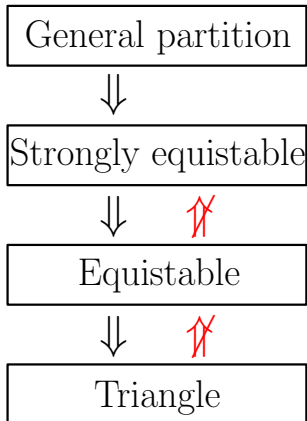
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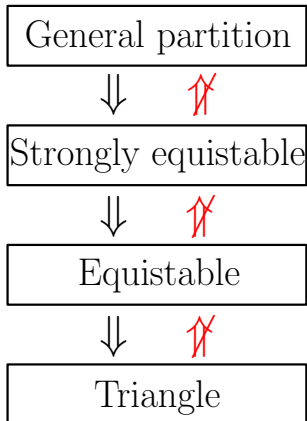
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Milanič, Trotignon, 2014

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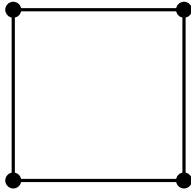
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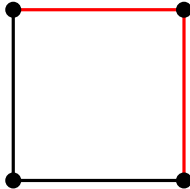
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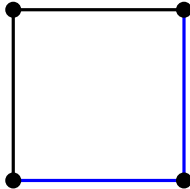


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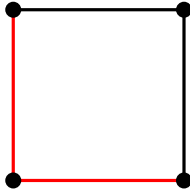




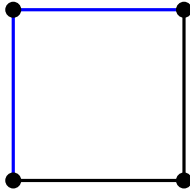
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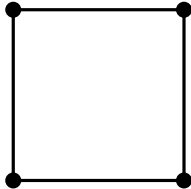
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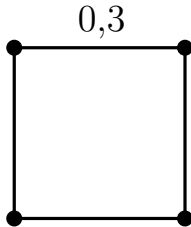
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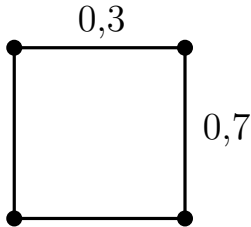
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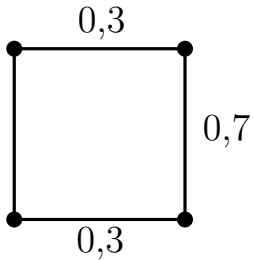
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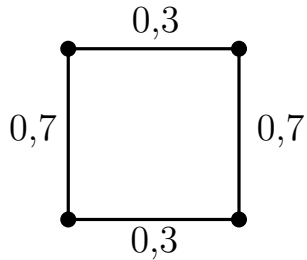
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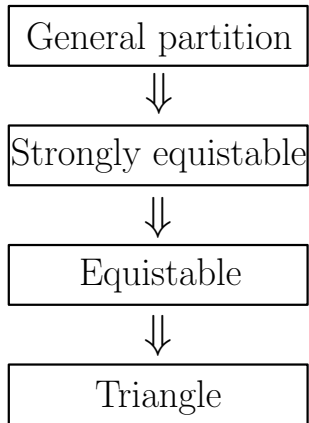
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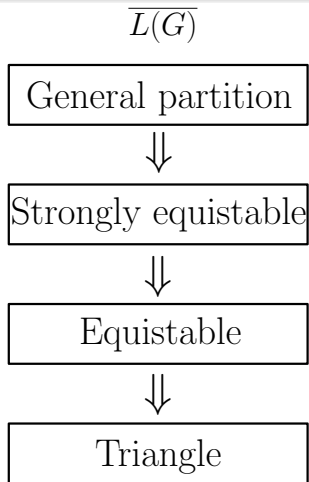
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# Equistarable graphs

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$G$

$\overline{L(G)}$

General partition



Strongly equistable



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Triangle

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triangle-free graph  $G$

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triangle-free graph  $G$

$\overline{L(G)}$

$\forall$  component is a star  
 or 2-internally extendable

$\iff$

General partition

$\Downarrow$

Strongly equistarable

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### Definition

A graph is 2- **internally extendable** if every 2-matching can be extended to a perfect internal matching.

triangle-free graph  $G$

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$\Downarrow$

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$P_5$ -constrained

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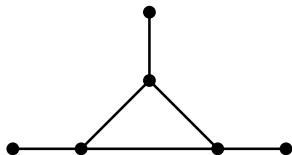
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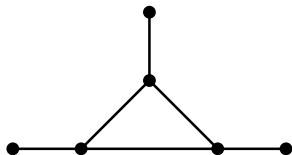
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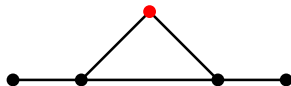
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not  $P_5$ -constrained

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# Bipartite graphs

A graph is **bipartite** if its vertex set can be partitioned into two stable sets.

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A graph is **bipartite** if its vertex set can be partitioned into two stable sets.

## Theorem

*For a bipartite graph  $G$  the following are equivalent:*

- (a) Every connected component of  $G$  is either a star or 2-internally extendable.*
- (b)  $G$  is strongly equistarable.*
- (c)  $G$  is equistarable.*

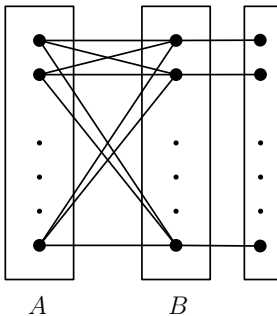
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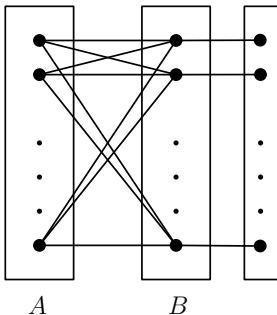
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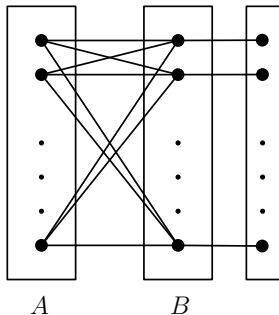
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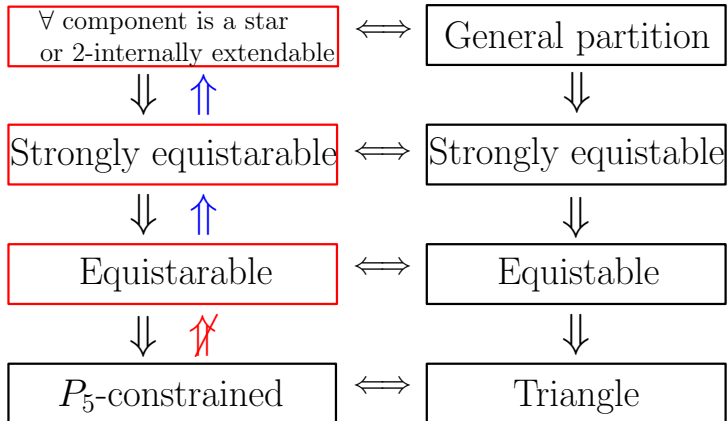
Every such graph with

$$3 \leq l \leq k+1$$

is not 2-internally extendable.

bipartite graphs

$\overline{L(G)}$



# Forests

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- (d)  $F$  is  $P_5$ -constrained.*

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Every tree  $T$  with  $|E(T)| \geq 1$  is almost 1-extendable.



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## Lemma

Every tree  $T$  with  $|E(T)| \geq 1$  is almost 1-extendable.

Let  $F$  be  $P_5$ -constrained. We can assume that  $F$  is connected and not a star.

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Every tree  $T$  with  $|E(T)| \geq 1$  is almost 1-extendable.

Let  $F$  be  $P_5$ -constrained. We can assume that  $F$  is connected and not a star.

Fix a 2-matching  $M = \{e, f\}$ , and consider the (unique) shortest path  $P$  between  $e$  and  $f$ .

We construct another matching  $M'$  by putting in for every vertex of  $P$ , not covered by  $M$ , an arbitrary edge incident with it and not in  $P$ .

(Since  $F$  is  $P_5$ -constrained, all the vertices of  $P$  have degree  $\geq 3$ .)

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Therefore, every connected component of  $F$  is either a star or 2-internally extendable.

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- Is every perfect equistable graph a general partition graph?

Thank you!